Arithmetic in Picard Groups and Application to Pairings 2

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Motivation

This talk is about fast arithmetic in divisor class groups of algebraic curves over finite fields for large genus.

What you do not get from this talk:

- Fast arithmetic for low genus curves optimised for use in a cryptographic system.

Some reasons why to consider this problem:

- General interest
- Useful for index calculus attacks, pairings.
Previous work

There is a long history of previous work on the theory and on algorithms for the

- Riemann-Roch problem
- arithmetic in class groups
- algebraic geometric codes
- integration of algebraic functions
- parametrisation of algebraic curves
- ...

Can roughly be divided into

- arithmetic methods (integral closures, ideals, ...)
- geometric methods (Brill-Noether method of adjoints, ...)
Previous work

Theory:

- Brill and Noether (1874, 1884),
- Dedekind and Weber (1882), F. K. Schmidt (1931).

**Geometric** and **arithmetic** algorithms for divisor class groups for $g \to \infty$:

<table>
<thead>
<tr>
<th>Year</th>
<th>Authors</th>
<th>Description</th>
<th>Complexity</th>
</tr>
</thead>
<tbody>
<tr>
<td>1987</td>
<td>Cantor</td>
<td>hyperelliptic divisor class group</td>
<td>$O(g^2)$</td>
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<tr>
<td>1993</td>
<td>Huang, Ierardi</td>
<td>RR problem + divisor class group for general plane curves</td>
<td>$O(n^6 h(D)^6)$</td>
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<tr>
<td>1994</td>
<td>Volcheck</td>
<td>divisor class group for general plane curves</td>
<td>$O(\max{n, g}^7)$</td>
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<tr>
<td>1998</td>
<td>Galbraith, Paulus, Smart</td>
<td>divisor class group for superelliptic curves</td>
<td>$O(n^4 g^4)$</td>
</tr>
<tr>
<td>1999</td>
<td>Arita</td>
<td>divisor class group for $C_{a,b}$ curves</td>
<td>$O(g^3)$</td>
</tr>
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</table>
Previous work

**Geometric** and **arithmetic** algorithms for divisor class groups for $g \to \infty$ (ctd):

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<th>Year</th>
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</tr>
</thead>
<tbody>
<tr>
<td>1999</td>
<td>Hess</td>
<td>RR problem and divclgrp for general curves</td>
<td>$O(g^2)$ for fixed $n$</td>
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<td>2001</td>
<td>Khuri-Makdisi</td>
<td>divclgrp for general curves with precomputation</td>
<td>$O^\sim(g^3)$</td>
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<td>2004</td>
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This and next slide $n = \min \{ [F : k(x)] \mid x \in F \text{ separating} \}$. 
Discussion

KM result:

• Links complexity of divclgrp to complexity of linear algebra over $k$ in dimension $O^\sim(g)$.
• Probably optimal in the general case ($n \gtrapprox g/2$).
• Fast linear algebra $O^\sim(g^\omega)$ with $\omega = 2.376$.

H result:

• Links complexity of divclgrp to complexity of polynomial arithmetic over $k$ in degree $O(g)$.
• Probably optimal under the assumption $n = O(1)$.
• Fast polynomial arithmetic $O^\sim(g)$.

This talk: Combine both running time characteristics towards $O^\sim(gn^{\omega-1})$ with $n = O(g)$. 
Divisor and ideal class groups

Let

- $x \in F$ be separating with $(x)_\infty = nP$ and $\deg(P) = 1$,
- $R = \text{Cl}(k[x], F)$.
- $n = O(g)$.

Then

- $R$ is a Dedekind domain.
- Ideals $I \neq \{0\}$ of $R$ are free $k[x]$-modules of rank $n$ and form a multiplicative monoid with cancellation law.
- $\text{Pic}(R) = (\text{group of fractional ideals})/(\text{group of principal ideals})$.
- $\text{Pic}(R) \cong \text{Pic}^0(F)$.
Arithmetic in the ideal class group

Represent ideal classes \([I]\) by integral ideals \(I\) of small „degree“.

Basic ideal operations for integral ideals \(I, J\):

- Simple multiplication: Compute \(zI\) for \(z \in J\).
- Integral division: Compute \(I/J\) for \(J \mid I\).

Degree reduction:

- \(R_z/I\) has small degree if \(z \in I\) has degree close to that of \(I\).
- Do not necessarily get unique reduction ... but yields bounded representation for \(I^{-1}\) and \(I\).

When writing \(z \in I\) we always mean to pick \(z\) of degree close to that of \(I\).
Arithmetic in the ideal class group

Arithmetic operations for \([I], [J] \in \text{Pic}(R)\):

- **Division:** \([I][J]^{-1} = [(zI)/J] \text{ for } z \in J\).
- **Inversion:** Use division with \([I] = [R]\).
- **Multiplication:** Use division and inversion.

Equality test for \([I], [J] \in \text{Pic}(R)\):

- **Let** \([K] = [I][J]^{-1}\).
- **Then** \([I] = [J] \text{ iff } K = Rz \text{ for some } z \in K \text{ of smallest degree.}\)

Use linear algebra over \(k[x]\)! 
Bases, matrices and degree function

Integral basis \( \omega_1, \ldots, \omega_n \in R \) of \( R \):

- \( \forall z \in R : \exists \) unique \( \lambda_i \in k[x] \) such that \( z = \sum_i \lambda_i \omega_i \).
- Multiplication table \( \lambda_{i,j,v} \in k[x] : \omega_i \omega_j = \sum_v \lambda_{i,j,v} \omega_v \).

Ideal basis \( \alpha_i \in I \) of ideal \( I \):

- \( \forall z \in I : \exists \) unique \( \lambda_i \in k[x] \) such that \( z = \sum_i \lambda_i \alpha_i \).
- Basis matrix \( M_I \in k[x]^{n \times n} : (\alpha_1, \ldots, \alpha_n) = (\omega_1, \ldots, \omega_n)M_I \).

Principal ideal \( I \):

- \( I = Rz \) for some \( z \in I \).
- Representation matrix \( M_z \in k[x]^{n \times n} : (z \omega_1, \ldots, z \omega_n) = (\omega_1, \ldots, \omega_n)M_z \).

Degree function:

- \( \deg^*(z) = -v_P(z) \) for \( z \in R \), \( \deg^*(I) = \deg(\det(M_I)) \).
- Have \( \deg^*(z) = \deg^*(Rz) \), \( \deg^*(x) = n \).
Bounded representations

Fix $\omega_1, \ldots, \omega_n$ with successively smallest $\deg^*$-values and let $d = g/n$.

**Theorem:**
1. There is a basis $\alpha_i$ of $I$ with $\deg^*(\alpha_i) \leq \deg^*(I) + 2g$ for all $i$.
   There is $\alpha_j$ with $\deg^*(\alpha_j) \leq \deg^*(I) + g$.
2. Elements of $\text{Pic}(R)$ can be represented by $I$ with $\deg^*(I) \leq g$.
3. $\deg(\lambda_{i,j,\nu}) \leq 4d$.
4. $\deg^*(\sum_i \lambda_i \omega_i) \leq w \Rightarrow \deg(\lambda_i) \leq w/n$ for all $i$.
5. $\deg^*(I) \leq w \Rightarrow$ there is a basis matrix $M_I$ with $\deg(M_I) \leq w/n$.

Represent elements of $\text{Pic}(R)$ by integral ideals of degree $\leq g$ with $n \times n$ basis matrices of degree $\leq 3d$. 
Linear algebra over polynomial rings

References: Storjohann, Villard, ...

Matrix multiplication in dimension $n$ and degree $d$:
- Time $O(d^2n^3)$.

Degree reduction (function field LLL, weak Popov form):
- Let $M = (v_1, \ldots, v_m) \in k[x]^{n \times m}$, $r$ be the rank of $M$, $d = \deg(M) = \max_i \deg(v_i)$ the maximum polynomial degree in $M$.
- $M$ is reduced iff $\deg(\sum_i \lambda_i v_i) = \max_i \deg(\lambda_i v_i)$ for all $\lambda_i \in k[x]$.
- $M$ can be transformed into reduced matrix by unimodular column operations in time $O(d^2nmr)$.

Kernel of $M$:
- Assume $M$ has a basis matrix $K$ for the $k[x]$-column kernel with $\deg(K) \leq d$ and that $m \geq n$.
- Then such a $K$ can be computed in time $O(d^2m^3)$.
Simple multiplication

Compute \( zI \) for \( z \in R \) with \( \deg^*(z) = O(g) \) and \( I \) integral ideal with \( \deg^*(I) = O(g) \).

Algorithm:

- Compute representation matrix \( M_z \) of \( z \) wrt \( \omega_i \).
  
  If \( z = \sum_i \mu_i \omega_i \) then \( z \omega_j = \sum_v (\sum_i \mu_i \lambda_{i,j,v}) \omega_v \).

- Multiply \( M_z \) and basis matrix of \( I \) to obtain a basis matrix of \( zI \).

Note \( \deg^*(zI) = \deg^*(z) + \deg^*(I) \).

Each step requires time \( O(d^2n^3) \).
Integral division

Let $I, J$ with $I \mid J$ and $\deg^*(J) = O(g)$. Compute $JI^{-1} = \{ z \in R \mid zI \subseteq J \}$.

- Let $I = \sum_{j=1}^{h} R\beta_j$ and $M_J$ be the basis matrix of $J$.
- For $z = \sum_i \lambda_i \omega_i$ and $\lambda = (\lambda_1, \ldots, \lambda_n)^t \in k[x]^n$:

$$z \in JI^{-1} \iff \exists v_i \in k[x]^n : \begin{pmatrix} M_{\beta_1} & M_J \\ \vdots & \vdots \\ M_{\beta_h} & M_J \end{pmatrix} \begin{pmatrix} \lambda \\ v_1 \\ \vdots \\ v_h \end{pmatrix} = 0.$$

Algorithm:

- Compute basis of kernel of big matrix, has rank $n$ and degree $O(d)$.
- Top $n \times n$ matrix yields basis matrix of $JI^{-1}$.

Required time $O(d^2(hn)^3)$.
Ideal basis reduction

Ideal basis reduction for $I$ with $\deg^*(I) = O(g)$:

- Let $d_i = \lceil \deg^*(\omega_i)/n \rceil$. Then $d_i = O(d)$.
- Let $M_I$ be a basis matrix of $I$ with $\deg(M_I) = O(d)$.

Algorithm:

- Multiply the $i$-th row of $M_I$ by $x^{d_i}$ for all $i$
- Apply the reduction algorithm.
- Divide the $i$-th row of the result by $x^{d_i}$ for all $i$.
- Denote the result by $M_I$.

The basis elements $\alpha_i$ then satisfy $\deg^*(\alpha_i) \leq \deg^*(I) + 2g$
and there is $\alpha_j$ with $\deg^*(\alpha_j) \leq \deg^*(I) + g$.

Required time $O(d^2n^3)$. 
Degree reduction

Degree reduction $I$ with $\text{deg}^*(I) = O(g)$:
- Let $M_I$ be a basis matrix of $I$ with $\text{deg}(M_I) = O(d)$.

**Algorithm:**
- Apply ideal basis reduction to $I$, yields new basis matrix $M_i$.
- Let $z = \alpha_j$ with $\text{deg}^*(\alpha_j) \leq \text{deg}^*(I) + g$.
- Apply simple multiplication and integral division to obtain $Rz/I$.

The ideal $Rz/I$ satisfies $\text{deg}^*(Rz/I) \leq g$. Apply procedure twice to reduce $I$.

Required time $O(d^2n^3)$. 
Principal ideal test

Principal ideal test for $I$ with $\deg^*(I) = O(g)$:

- $\deg^*(z) \geq \deg^*(I)$ for all $z \in I$,
- $I = Rz$ iff $z \in I$ and $\deg^*(z) = \deg^*(I)$.

- Let $\alpha_i$ be a reduced ideal basis.
- The ideal basis reduction also yields integers $e_1 \leq \cdots \leq e_n$ with
  $$L(I, r) = \{z \in I \mid \deg^*(z) \leq rn\} = \{\sum_i \lambda_i \alpha_i \mid \deg(\lambda_i) \leq -e_i + r\}$$
  for all $r \in \mathbb{Z}$.
- If $z \in R$ such that $\deg^*(zI) = rn$, then $zI$ principal iff $L(zI, r) \neq 0$.

Algorithm:

- Compute $z \in R$ such that $\deg^*(zI) = rn$ and $\deg^*(z) = O(g)$.
  (Precomputation for all possible degrees of $I$).
- Using ideal basis reduction on $zI$ check $L(zI, r) \neq 0$.

Required time $O(d^2n^3)$. 
Ideal generating sets

Time for integral division is $O(d^2(hn)^3)$.

Let $I$ be an ideal with $\deg^*(I) = O(g)$ and reduced basis $\alpha_i$.
Let $h = \max\{\log_q(g), 2\}$.

Proposition (KM):
- A random choice of $h$ elements $\beta_j$ of $\sum_{i=1}^{n} k\alpha_i$ is a generating system for $I$ with probability $\geq 1/2$.

Algorithm for integral division:
- Choose $h$ random such $\beta_j$ (for $n = O(1)$ we can take the $\alpha_i$).
- Compute reduced basis of $J/\sum_j R\beta_j$.
- If $\deg^*(J/\sum_j R\beta_j) \neq \deg^*(J) - \deg^*(I)$ then repeat.

Required expected time $O^\sim(d^2 n^3)$. 
Multiplication table speed up

Time for representation matrix computation $O(d^2n^3)$.

Use FFT inspired/based technique:

- Work in $R/Rx^{O(d)}$, can lift elements of small degree back to $R$.
- Assume there is $y \in R$ such that \(\gcd((R : k[x,y]), x) = 1\).
- Then $R/x^d \cong k[x,y]/x^d$, are free $k[x]/x^d$-modules.
- Isomorphism evaluated by multiplication of unimodular $n \times n$ matrix of degree $\leq d$.
- Multiplication in $k[x,y]/x^d$ in time $O(d^2n^2)$ or $O^\prec(dn)$. 
Pairing computation

Need to evaluate the functions $z \in R$ occurring in simple multiplication and principal ideal test:

- Compute norm of $z$ in $R/J$ over $k$.
- If zero then choose different equivalent $J$.
- No inversions necessary.
- Maybe can use $J$ for FFT inspired technique simultaneously (counted before as precomputation).
- $O(\log_2(r))$ evaluations in total.

Runtime not evaluated yet ...
Conclusion

The overall running time for arithmetic in $\text{Pic}^0(F)$ is

$$O^\sim(d^2n^3) = O^\sim(g^2n)$$

where $dn = g$.

- For $d = O(1)$ we obtain $O^\sim(g^3)$ (KM).
- For $n = O(1)$ we obtain $O(g^2)$ (H).
- For $C_{a,b}$ curves we obtain $O^\sim(g^{5/2})$.

The running time linkable to linear algebra over polynomial rings with $nd = g$, should result in $O^\sim(dn^\omega) = O^\sim(gn^{\omega-1})$.

An $n = O(1)$ and time $O(g^2)$ or $O^\sim(g)$ implementation is available in the computer algebra systems Kash and Magma.

Pairings should be similar ... runtime $O^\sim(\log_2(r)gn^{\omega-1})$. 

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